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INTENSIONALLY COMPLEMENTED DISTRIBUTIVE LATTICES*

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It is well known that plausible semantic completeness and consistency conditions for ordinary classical logic are possible in terms of boolean algebras ([11]), and that something similar is possible for both intuitionistic logic and for modal logics via topological structures ([13]). An analogous problem is the provision of a suitable semantics for the intensional system E of entailment ([1], [2], [3], [6], [7]). Although the general problem remains unsolved, it has nevertheless been possible to provide solutions for at least certain fragments of E ([4], [5], [8]). These partial solutions to the semantic problem implicitly utilize a certain type of lattice, as follows: an ordered quadruple $\langle A, \leq, N, T \rangle$ (where A is a set, \leq a relation on A , N a function on A , and T a subset of A) is said to be an *intensionally complemented distributive lattice with truth-filter* (icdl w/t-f) provided

- (DL) A is a distributive lattice under \leq (we use \cap and \cup for the lattice meet and join respectively, as in [10]);
- (N₁) $NNa = a$, all $a \in A$;
- (N₂) if $a \leq b$, then $Nb \leq Na$, all $a, b \in A$;
- (T) T is a *truth-filter*; i. e., T satisfies these three conditions: (i) T is a *filter* ([15]): $(a \cap b) \in T$ if both $a \in T$ and $b \in T$, and $b \in T$ if both $a \leq b$ and $a \in T$; (ii) T is *consistent*: there is no a such that both $a \in T$ and $Na \in T$; and (iii) T is *exhaustive*: for every $a \in A$, either $a \in T$ or $Na \in T$.

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In the semantic applications to intensional logic, A is thought of as a set of propositions, \leq as a relation of entailment, the lattice meet and join as propositional conjunction and disjunction respectively, and N as propositional negation. T is then thought of as representing the «true» propositions in the set A .

In these applications the truth-filter T is employed explicitly; however, certain of the semantic results can be seen to hold also where the notion of an *icdl w/t*— f is replaced by the more general notion of an *intensionally complemented distributive lattice*, i.e., a triple $\langle A, \leq, N \rangle$ such that for some T , $\langle A, \leq, N, T \rangle$ is an *icdl w/t*— f .

Separate study has also been given structures satisfying (DL), (N_1) , and (N_2) , which are *de Morgan Lattices* in the sense of [14]. (They have also been called *quasi-boolean algebras* in [9], and *distributive i-lattices* in [12].) Comparison of the two notions naturally leads to the request for a necessary and sufficient condition that a de MORGAN lattice be an *icdl*. It then turns out that a quite simple condition will do, namely, that N have no fixed point :

$$(N_3) \quad a \neq Na, \quad \text{all } a \in A.$$

Two remarks are in order before turning to a proof of the equivalence of (N_3) and (T). (i) The equivalence gains interest from the extremely high intuitive plausibility of (N_3) and (T) as conditions on propositional negation. Indeed, it is hard to think of anything more absurd than supposing a proposition could be its own negation, and it is hardly less absurd to suppose that it should be impossible to exactly divide the propositions into the False and the True, a proposition being true iff its negation is false, and a conjunction being true iff each conjunct is true. (ii) In [14] a de MORGAN lattice, A , is said to be *essentially irregular* if there is no set D such that $\langle A, D \rangle$ is a regular characteristic matrix for the classical propositional calculus (see references in [14]). It is a corollary of the chief result of this paper, together with Theorem 4 of [14], that a necessary and sufficient condition for essential irregularity is the presence of some fixed point for N .

Since the equivalence of (N_3) and (T) is obvious in one direction, it suffices to prove the following

THEOREM. *If A is a de Morgan Lattice such that N has no fixed point, then some filter T of A is both consistent and exhaustive.*

Were A a boolean algebra, the theorem would amount to the existence of maximal proper filters; but where the complementation is intensional, not every proper filter is consistent nor is every exhaustive filter maximal.

Turning now to a proof of the theorem, we first remark that all the various forms of de MORGAN'S law hold, a fact to which we refer below by means of «(deM)». Now let T_0 be the set of all elements q of A such that $q = (q \cup Nq)$, and let E be the set of all consistent filters in A containing T_0 . We first show E non-empty by showing consistent the filter $F(T_0)$ generated by T_0 . For suppose otherwise; then for some $p_1, \dots, p_n \in T_0$ and some $b \in A$, we would have, for $p = p_1 \cap \dots \cap p_n$, both $p \leq b$ (hence $Nb \leq Np$) and $p \leq Nb$, so that

$$(1) \quad Np = p \cup Np.$$

Now define a_n inductively as follows:

$$(2) \quad a_1 = p_1; \quad a_{i+1} = p_{i+1} \cap (Np_{i+1} \cup a_i).$$

We first prove inductively that

$$(3) \quad a_i = (p_1 \cap \dots \cap p_i) \cup N a_i, \quad \text{all } i \quad (1 \leq i \leq n).$$

$a_1 = p_1 \cup N a_1$ is immediate, and we now assume for induction that $a_i = (p_1 \cap \dots \cap p_i) \cup N a_i$. Since $p_{i+1} \in T_0$, we have $p_{i+1} = (p_{i+1} \cup N p_{i+1})$, hence by (DL), $p_{i+1} \cap (N p_{i+1} \cup a_i) = (p_{i+1} \cap a_i) \cup N p_{i+1}$; so that by (2),

$$(4) \quad a_{i+1} = (p_{i+1} \cap a_i) \cup N p_{i+1}.$$

By the hypothesis of the induction, (4) yields

$$\begin{aligned} (5) \quad a_{i+1} &= (p_{i+1} \cap ((p_1 \cap \dots \cap p_i) \cup N a_i)) \cup N p_{i+1} \\ &= (p_1 \cap \dots \cap p_{i+1}) \cup (N p_{i+1} \cup (p_{i+1} \cap N a_i)) && \text{by (DL)} \\ &= (p_1 \cap \dots \cap p_{i+1}) \cup N (p_{i+1} \cap (N p_{i+1} \cup a_i)) && \text{by (deM)} \\ &= (p_1 \cap \dots \cap p_{i+1}) \cup N a_{i+1} && \text{by (2)} \end{aligned}$$

Hence (3) holds, which can be used to show that $a_n = N a_n$ as follows:

$$\begin{aligned}
(6) \quad a_n &= \not{p} \cup N a_n && \text{by (3)} \\
&= \not{p} \cup N(\not{p} \cup N a_n) && \text{by (3)} \\
&= \not{p} \cup (N \not{p} \cap a_n) && \text{by (deM)} \\
&= (\not{p} \cup N \not{p}) \cap (\not{p} \cup a_n) && \text{by (DL)} \\
&= N \not{p} \cap (\not{p} \cup a_n) && \text{by (1)} \\
&= N \not{p} \cap (\not{p} \cup (\not{p} \cup N a_n)) && \text{by (3)} \\
&= N \not{p} \cap (\not{p} \cup N a_n) && \text{by (DL)} \\
&= N \not{p} \cap a_n && \text{by (3)} \\
&= N(\not{p} \cup N a_n) && \text{by (deM)} \\
&= N a_n && \text{by (3)}
\end{aligned}$$

But this contradicts (N_3) , so that $F(T_0)$ must after all be consistent, and the set E of all consistent filters in A containing T_0 must be non-empty. Since it is easy to see that if a subset G of E is simply ordered by \subseteq , then $\bigcup G \in E$, it follows by a form of ZORN'S Lemma that E has a maximal element T . T is obviously consistent; we proceed to show that it is exhaustive as well, which will complete the proof of the theorem.

We observe to begin with that

$$(7) \quad (a \cup N a) \in T, \quad \text{all } a \in A,$$

since $(a \cup N a) = ((a \cup N a) \cup N(a \cup N a))$ by (deM) and (DL); so that $(a \cup N a) \in T_0$. Furthermore, since T is a filter, it follows from (7) that $((c \cup N c) \cup (d \cap N d)) \in T$ for all $c, d \in A$; hence, since T is consistent,

$$(8) \quad N((c \cup N c) \cap (d \cup N d)) \notin T, \quad \text{all } c, d \in A.$$

Now suppose for *reductio* that T is not exhaustive, so that for some b , neither $b \in T$ nor $Nb \in T$. Then the filters $F(T, b)$ and $F(T, Nb)$ generated by T with b and T with Nb respectively each have T as a proper subset; and so since T is maximal in E , neither $F(T, b)$ nor $F(T, Nb)$ can be consistent. Consequently, for some c , both c and Nc are in $F(T, b)$, and for some d , both d and Nd are in $F(T, Nb)$; so since $F(T, b)$ is the filter generated by T with b and since $F(T, Nb)$ is the filter generated by T with Nb , it must be that for some $t_1, t_2, t_3, t_4 \in T$,

$$(9) \quad t_1 \cap b \leq c, \quad t_2 \cap b \leq Nc, \quad t_3 \cap Nb \leq d, \quad \text{and } t_4 \cap Nb \leq Nd,$$

which leads by (D L) to

$$(10) \quad t_1 \cap t_2 \cap t_3 \cap t_4 \cap (b \cup N b) \leq (N c \cap c) \cup (N d \cap d).$$

Since $t_1, t_2, t_3, t_4 \in T$, since also $(b \cup N b) \in T$ by (7), the fact that T is a filter implies that $(t_1 \cap t_2 \cap t_3 \cap t_4 \cap (b \cup N b)) \in T$, and accordingly, with (10), that $((N c \cap c) \cup (N d \cap d)) \in T$. So by (deM)

$$(11) \quad N((c \cup N c) \cap (d \cup N d)) \in T.$$

But (11) contradicts (8), rendering absurd the supposition that T fails to be exhaustive.

This completes the proof of the theorem. We remark that it cannot be generalized to *complete* distributive lattices (where meets and joins of arbitrary subsets always exist) satisfying $(N_1) - (N_3)$ and *complete* filters (filters closed under the meet of arbitrary subsets). The following is an example of a complete distributive lattice satisfying $(N_1) - (N_3)$ which contains no complete truth-filter: let $\langle A, \leq, N \rangle$ be such that (i) the elements of A are $(1, \infty)$ and $(-1, -\infty)$ together with all ordered pairs (x, y) where $x = \pm 1$ and y is an integer; (ii) $(x_1, y_1) \leq (x_2, y_2)$ iff $x_1 \leq x_2$ and $y_1 \leq y_2$; and (iii) $N(x, y) = (-x, -y)$. It is easily shown that $\langle A, \leq, N \rangle$ is a complete distributive lattice satisfying $(N_1) - (N_3)$. Suppose now there is a complete, consistent, and exhaustive filter T in A . Then since $(-1, 0) \leq N(-1, 0)$, we shall have to have $(-1, 0) \notin T$, on pain of inconsistency of T ; hence $N(-1, 0) \in T$, since T is exhaustive. However, for all $x \geq 0$, $(-1, x) \cap N(-1, 0) = (-1, 0)$, so that the filterhood of T implies that $(-1, x) \notin T$ for all such x . But T is exhaustive, so the set X of all pairs $N(-1, x) [= (1, -x)]$, for $x \geq 0$, is a subset of T ; and since T is a complete filter, the meet of X is in T . But $(-1, -\infty)$ is evidently the meet of X so that $(-1, -\infty) \notin T$ —from which it follows immediately that T is inconsistent, contrary to assumption.

We may also observe that the theorem cannot be strengthened to read «every consistent filter is a subset of some exhaustive consistent filter»; a counter-example is provided by $\langle B, \leq, N \rangle$, where $B = \{(-1, -\infty), (1, \infty), (-1, 0), (1, 0), (-1, 1), (1, -1)\}$ and where \leq and N are as in the previous example: the filter $\{(-1, 1), (1, \infty)\}$ is consistent, but not extendible to an exhaustive consistent filter.

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